

Challenge. Without using Zorn, devise a spell for switching on Zorn: construct a Grothendieck topos with surjective geometric morphism to Set in which Zorn holds.



Pointfreely embracing infinite sequences as useful fictions

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What is a function in constructive mathematics?

Function types are brittle.

(naive) statement	cl.	sp. in Eff	
0 Every function is computable.	✗	?	✓
1 Every function has a zero or not.	✓	?	✗
2 Every path leaves the Kleene tree.	✗	?	✓
3 There is no surjection $\mathbb{N} \rightarrow \mathbb{R}$.	✓	?	✓

False in parametrized realizability.

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3 There is no surjection $\mathbb{N} \rightarrow \mathbb{R}$.	✓	? ✓	False in parametrized realizability.

With **adverse effects** throughout mathematics:

4 Dickson's lemma.	✓	? ✗	"Products of well quasi-orders are well."
5 Products of streamless sets are streamless.	✓	? ✗	With LEM, true in Gr. toposes with enough points.
6 Products of wf. relations are wf.	✓	?	
7 Minimum principle for unif. cont. functions.	✓	? ✗	"For $f : [0, 1] \rightarrow \mathbb{R}_{>0}$, there is $c > 0$ with $f \geq c$."
8 Every field has an algebraic closure.	✓	? ✗	Implies a form of choice.
9 Hilbert's basis theorem.	✓	?	"If A is Noetherian, so is $A[X]$."
10 Locally Noetherian implies Noetherian.	✓	?	"If all $A[f_i^{-1}]$ are Noetherian, so is A ."
11 Krull's lemma.	✓	?	"If $x \in \bigcap_{\mathfrak{p}} \mathfrak{p}$, then x is nilpotent."
12 Fundamental theorem of Galois theory.	✓	?	"If $x \in L^{\text{Aut}_K(L)}$, then $x \in K$."

Harnessing formal topology

Let X be a set.

Irrespective of **function types**, there is a **robust formal space**^{*} $X^{\mathbb{N}}$ with

- **basic opens:** $[x_0, \dots, x_{n-1}]$, pictured as “ $\{f : \mathbb{N} \rightarrow X \mid f(0) = x_0, \dots, f(n-1) = x_{n-1}\}$ ”
- **coverings inductively generated by:** $[x_0, \dots, x_{n-1}] \triangleleft \{[x_0, \dots, x_{n-1}, a] \mid a \in X\}$

^{*} distributive lattice, site, locale, topos, arithmetic universe, ...

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We can then explore:

$$\text{(absolute) points of } X^{\mathbb{N}} \quad \xleftrightarrow{1:1} \quad \text{functions } \mathbb{N} \rightarrow X$$

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For suitable spaces E , E -valued points of $X^{\mathbb{N}}$ are for instance:

- Partially defined functions $f : \mathbb{N} \multimap X$ such that for all n , the value $f(n)$ is *not not* defined
- Total multivalued functions $f : \mathbb{N} \rightrightarrows X$ (approximately)

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Dialogues with oracles

For a predicate P on $\text{List}(X)$, we inductively define $P \mid \sigma$ (“ P bars the list σ ”) by the following two clauses.

- 1 “now”: If $P \sigma$, then $P \mid \sigma$.
- 2 “later”: If $P \mid \sigma a$ for all $a \in X$, then $P \mid \sigma$.

The assertion $P \mid \sigma$ expresses in a positive way that **no matter how** the finite approximation σ evolves over time to a **better approximation** τ , **eventually** $P \tau$ will hold.

A witness of $P \mid \sigma$ is a **dialogue tree** recording **queries** to an oracle of type $\mathbb{N} \rightarrow X$, eventually culminating with witnesses of P at the leaves.

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A brittle definition

A set X is **streamless** iff
for every function $f : \mathbb{N} \rightarrow X$,
there are numbers $i < j$ such that $f(i) = f(j)$.

Stability of streamless sets under products
depends on foundational assumptions.

A robust definition

A set X is **Noetherian** iff $R \mid []$,
where $R [x_0, \dots, x_{n-1}]$ iff
there are numbers $i < j$ such that $x_i = x_j$.

Stability of Noetherian sets under products is
provable in spartan mathematics.

Elephant functions

Def. Let X and Y be spaces. An **elephant function** from X to Y consists of

■ for each space E , a function $f_E : \text{Pt}_E(X) \rightarrow \text{Pt}_E(Y)$ between the sets of E -valued points such that

$$\text{Pt}_E(X) \xrightarrow{f_E} \text{Pt}_E(Y)$$

$$\circ g \downarrow \qquad \qquad \qquad \downarrow \circ g \quad \text{commutes.}$$

$$\text{Pt}_F(X) \xrightarrow{f_F} \text{Pt}_F(Y)$$

So a natural transformation $\mathcal{Y}X \rightarrow \mathcal{Y}Y$, so equivalently a morphism $X \rightarrow Y$ of spaces.

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■ for each morphism $g : F \rightarrow E$, the diagram

So a natural transformation $\mathcal{L}X \rightarrow \mathcal{L}Y$, so equivalently a morphism $X \rightarrow Y$ of spaces.

Prop. For a function $f : (\mathbb{N} \rightarrow \mathbb{2}) \rightarrow \mathbb{N}$, TFAE:

- 1 f can be extended to an elephant function from $\mathbb{2}^{\mathbb{N}}$ to \mathbb{N} .
- 2 f is uniformly continuous.
- 3 f is eloquent (i.e. given by a dialogue).

Prop. For a function $f : [0, 1] \rightarrow \mathbb{R}$, TFAE:

- 1 f can be extended to an elephant function from $[0, 1]$ to \mathbb{R} .
- 2 f is uniformly continuous.
- 3 f is eloquent (in a suitable real-number-oracle querying sense).

Def. A statement φ holds ...

- **everywhere** ($\Box\varphi$) iff it holds **over every space**.
- **somewhere** ($\Diamond\varphi$) iff it holds **over some positive space**.
(A space is *positive* iff every open covering of its top open is inhabited).
- **proximally** ($\Diamond\varphi$) iff it holds **over some positive overt space**.

Modal language for embracing points

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Equivalently: $\Box(X \text{ is streamless})$.

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A brittle definition

A relation R is **well-founded _{∞}** iff
for every decreasing sequence $f : \mathbb{N} \rightarrow X, \perp$.


A robust definition

A relation R is **well-founded_{ind}** iff $\text{Acc}(R)$.

Equivalently: $\Box(R \text{ is well-founded}_{\infty})$.

- 1 $(\Diamond\varphi) \iff \varphi$, if φ is a geometric implication (“ $\forall \dots \forall(\% \Rightarrow \%)$ ”, with no \forall nor \Rightarrow in $\%$). ⚙
- 2 $(\Diamond\varphi) \iff \varphi$, if φ is a bounded first-order statement. ⚙

Harmonized results

statement	cl.	sp.	in Eff
0 <i>Everywhere</i> , every function is computable.	✗	✗	✗
1 <i>Everywhere</i> , every function has a zero or not.	✗	✗	✗
2 <i>Everywhere</i> , every path leaves the Kleene tree.	✗	✗	✗
3 <i>Everywhere</i> , there is no surjection $\mathbb{N} \rightarrow \mathbb{R}$.	✗	✗	✗
4 Products of well _{ind} quasi-orders are well _{ind} .	✓	✓	✓
5 Products of Noetherian sets are Noetherian.	✓	✓	✓
6 Products of well-founded _{ind} relations are well-founded _{ind} .	✓	✓	✓
7 If <i>everywhere</i> $f(x) > 0$ for all x , then there is $c > 0$ with $f \geq c$.	✓	✓	✓
8 Every field has an algebraic closure <i>somewhere</i> .	✓	✓	✓
9 If A is Noetherian _{ind} , so is $A[X]$.	✓	✓	✓
10 Locally Noetherian _{ind} implies Noetherian _{ind} .	✓	✓	✓
11 If <i>everywhere</i> x is included in every prime ideal, then x is nilpotent.	✓	✓	✓
12 If <i>everywhere</i> $\sigma(x) = x$ holds for all $\sigma \in \text{Aut}_K(L)$, then $x \in K$.	✓	✓	✓
13 For every inhabited set X , <i>proximally</i> there is a surjection $\mathbb{N} \twoheadrightarrow X$.	✓	✓	✓ 
14 <i>Somewhere</i> , LEM holds.	✓	✓	✓
15 Barr's theorem: $\text{Zorn} \Rightarrow \square \diamond \text{AC}$.	✓	✓	✓

Taming the infinite

Def. Let A be a ring.

- An infinite sequence x_0, x_1, \dots *stalls* iff there is a number n with $x_n \in (x_0, \dots, x_{n-1})$.
- A finite list $[x_0, \dots, x_{m-1}]$ *stalls* iff there is a number $n < m$ with $x_n \in (x_0, \dots, x_{n-1})$.

Notation: “ $\text{Stalls}_A [x_0, \dots, x_{m-1}]$ ”

Def. A ring A is ...

- Noetherian_∞ iff every infinite sequence of elements of A stalls.
- $\text{Noetherian}_{\text{ind}}$ iff $\text{Stalls}_A \mid []$.
Equivalently: $\square(A \text{ is } \text{Noetherian}_\infty)$.

Thm. (classical) Let A be a Noetherian_∞ ring. Then $A[X]$ is Noetherian_∞ as well.

Proof. See [Hilbert 1890]. *This is not mathematics, it is theology!* □

Thm. Let A be a $\text{Noetherian}_{\text{ind}}$ ring. Then $A[X]$ is $\text{Noetherian}_{\text{ind}}$ as well.

Proof. Let an arbitrary sequence of elements of $A[X]$ be given (over an arbitrary space). It suffices to verify that it stalls *somewhere*. *Somewhere*, LEM holds, and—crucially—*there* A is still Noetherian_∞ . So *there* the sequence stalls by the classical proof. □

Answering a question by Berardi–Buriola–Schuster

Def. Let (X, \leq) be a quasi-order.

- An infinite sequence x_0, x_1, \dots is *good* iff there are numbers $i < j$ with $x_i \leq x_j$.
- A finite list $[x_0, \dots, x_{m-1}]$ is *good* iff there are numbers $i < j < m$ with $x_i \leq x_j$.
Notation: “Good $[x_0, \dots, x_{m-1}]$ ”

Def. A quasi-order X is ...

- well_∞ iff every infinite sequence is good.
- well_{ind} iff Good $| []$ (equivalently: *everywhere*, X is well_∞).
- $\text{well}_{\text{impl}}$ iff for every monotone predicate B of finite lists, $B \upharpoonright_{\text{incr}} []$ implies $B \upharpoonright []$ (equivalently: if, everywhere, every *increasing* infinite sequence $\alpha : \mathbb{N} \rightarrow X$ has a finite prefix validating B , then, everywhere, every infinite sequence does).

Thm. Let X be a well_{ind} quasi-order. Then X is $\text{well}_{\text{impl}}$.

Proof. Let B be a monotone predicate. Let x_0, x_1, \dots be an infinite sequence (over an arbitrary space). It suffices that it has a finite B -prefix *somewhere*. *Somewhere*, LEM holds. By a classical lemma, *there* we have an increasing infinite subsequence x_{i_0}, x_{i_1}, \dots . By assumption, *there* we have a finite prefix with $B[x_{i_0}, \dots, x_{i_n}]$. As B is monotone, *there* we also have $B[x_0, x_1, \dots, x_{i_n}]$. □

Mismatches

Let A be a ring.

1 Let M be an A -module.

The (robust) statement “*everywhere*, for every prime filter \mathfrak{p} the stalk $M_{\mathfrak{p}}$ is finitely generated” is classically **not equivalent** (but stronger) than its single-worlded unmodal counterpart “for every absolute prime filter \mathfrak{p} , the stalk $M_{\mathfrak{p}}$ is finitely generated”.

2 Let $f : A \rightarrow B$ be a homomorphism.

The status of the (robust) statement “*everywhere*, the preimage of every maximal ideal of B is maximal in A ” is **currently unclear**.

It might unfortunately be strictly stronger, classically, than “the preimage of every maximal ideal of B is maximal in A ”.

(Compare with Ryota Kuroki’s pioneering work on the Nullstellensatz and Jacobson rings in constructive algebra.)

Embracing implication as a useful fiction

(Bonus slide for people who are intrigued by arithmetic universes.)

Let \mathbb{T} be a geometric theory. Let $U_{\mathbb{T}}$ be the generic \mathbb{T} -model in $\text{Sh}(\mathcal{C}_{\mathbb{T}})$.

Let $\mathbb{T}/U_{\mathbb{T}}$ be the theory (in $\text{Sh}(\mathcal{C}_{\mathbb{T}})$) of $U_{\mathbb{T}}$ -algebras.

Let A and B be geometric formulas in the language of $\mathbb{T}/U_{\mathbb{T}}$ in a context \vec{x} . Then the assertion

$$“\forall \vec{x}. (A \Rightarrow B)”$$

makes use of implication and can therefore only be represented in $\text{Sh}(\mathcal{C}_{\mathbb{T}})$ if Set supports implication. (It references the **just one point** of the classifying space of $\mathbb{T}/U_{\mathbb{T}}$, the initial one.)

In contrast, the assertion

$$“\mathbb{T}/U_{\mathbb{T}} \vdash (\forall \vec{x}. (A \Rightarrow B))”$$

is positive (a geometric formula) and hence available even if implication is not. (It can be understood as referencing the **generic point** of the classifying space of $\mathbb{T}/U_{\mathbb{T}}$.)

Curiously, these two assertions are equivalent (in a metatheory with implication).

So use the latter as a proxy for the former!