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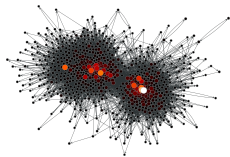
$$\begin{pmatrix} a_{11}(t) & \cdots & a_{1n}(t) \\ \vdots & & \vdots \\ a_{n1}(t) & \cdots & a_{nn}(t) \end{pmatrix}$$

There for every parameter value  $t$ , classically there is

- ▶ a full list of eigenvalues  $\lambda_1(t), \dots, \lambda_n(t)$  and
- ▶ an eigenvector basis  $(v_1(t), \dots, v_n(t))$ .



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How about the  $v_i$ ?



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Can locally the functions  $\lambda_i$  be chosen to be continuous? **Yes.**  
How about the  $v_i$ ? **No.**

— an invitation —

# New modal operators for constructive mathematics

*Type Theory, Constructive Mathematics and Geometric Logic*

CIRM

May 2nd, 2023

Ingo Blechschmidt

j.w.w. Alexander Oldenziel

# Questions

- 1 Why has the inductive revolution been so powerful?
- 2 Why do proofs using Zorn's maximal ideals work so well in constructive algebra?
- 3 Why are elements of  $\bigcap_{\mathfrak{p}} \mathfrak{p}$  not necessarily nilpotent?
- 4 How can we extract computational content from classical proofs?

# Infinite data

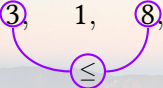
7, 4, 3, ...

# Infinite data

7, 4, 3, 1, ...

# Infinite data

7, 4, 3, 1, 8, ...



A diagram illustrating a sequence of numbers: 7, 4, 3, 1, 8, ... The numbers 3 and 8 are circled in purple. A purple curved line connects the circle around 3 to a circle around the less-than-or-equal-to symbol ( $\leq$ ), which is then connected by another purple curved line to the circle around 8. This indicates a comparison between the values 3 and 8.

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**Thm.** Every sequence  $\alpha : \mathbb{N} \rightarrow \mathbb{N}$  is **good** in that there exist  $i < j$  with  $\alpha(i) \leq \alpha(j)$ .



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*Proof.* (offensive?) By **LEM**, there is a minimum  $\alpha(i)$ . Set  $j := i + 1$ . □

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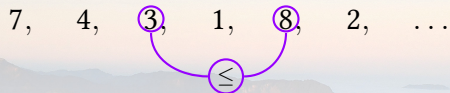
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**Examples.**  $(\mathbb{N}, \leq)$ ,  $X \times Y$ ,  $X^*$ ,  $\text{Tree}(X)$ .

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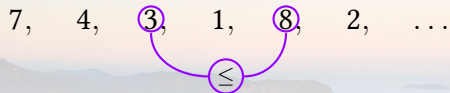
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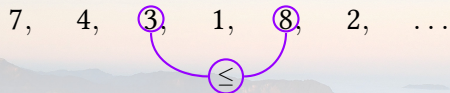
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# Computational content from classical proofs

**Def.** A transitive relation ( $<$ ) on a set  $X$  is ...

1 **well-founded\*** iff there is no **infinite chain**  $x_0 > x_1 > \dots$ ,

2 **well-founded** iff for every  $x \in X$ ,  $\text{Acc}(x)$ ,

where  $\text{Acc}$  is inductively defined by:

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**Prop.** Let  $(X, \leq)$  be preorder. Let “ $x < y$ ” mean  $x \leq y \wedge \neg(y \leq x)$ .  
Then: If  $X$  is  $\text{well}^*$ , then  $(<)$  is  $\text{well-founded}^*$ .

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Can we extract a constructive proof that well preorders are well-founded?



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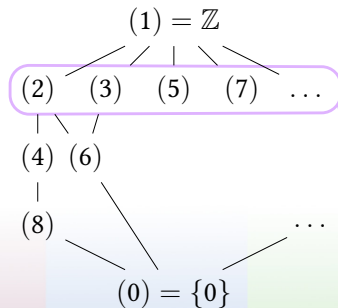
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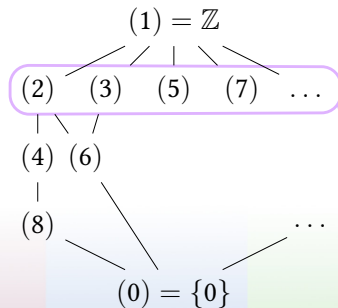


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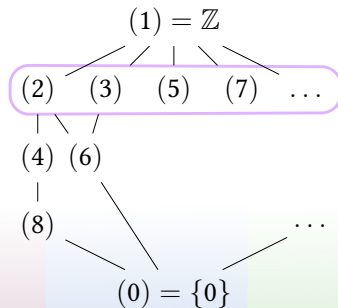
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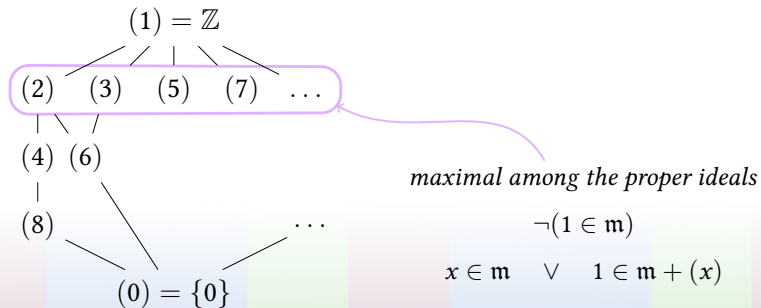
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$$\mathfrak{m}_0 := \{0\}, \quad \mathfrak{m}_{n+1} := \begin{cases} \mathfrak{m}_n + (x_n), & \text{if } 1 \notin \mathfrak{m}_n + (x_n), \\ \mathfrak{m}_n, & \text{else.} \end{cases}$$

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4 In the general case: **No**,  
but *first-order consequences* of the existence of a maximal ideal **do hold**.

# Questions

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# Constructive forcing (= Grothendieck toposes)

Let  $L$  be a **forcing notion**, a preorder equipped with a **covering system**.<sup>1</sup> **Filters**  $F \subseteq L$  are subsets which are upward-closed, downward-directed and split the covering system.<sup>2</sup>

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<sup>1</sup>A covering system consists of a set  $\text{Cov}(\sigma) \subseteq P(\downarrow\sigma)$  of *coverings* for each element  $\sigma \in L$  subject only to the following simulation condition: If  $\tau \preceq \sigma$  and  $R \in \text{Cov}(\sigma)$ , there should be a covering  $S \in \text{Cov}(\tau)$  such that  $S \subseteq \downarrow R$ .

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**Def.** Given a monotone predicate  $P$  on  $L$ , inductively define:

$$\frac{P\sigma}{P \mid \sigma} \quad \frac{\forall (\tau \in R). P \mid \tau}{P \mid \sigma} \quad (R \in \text{Cov}(\sigma))$$

We use quantifier-like notation: “ $\nabla \sigma. P\sigma$ ” means  $P \mid \sigma$ .

# Constructive forcing (= Grothendieck toposes)

A forcing notion is a template for a **forcing extension**  $V^\nabla$  of the base universe  $V$ :

*When we say that a statement **holds in**  $V^\nabla$ ,  
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- 4 Let  $X$  be a preorder. Then:
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# The modal multiverse

In general, “ $\varphi$  holds in  $V^\nabla$ ” and “ $\varphi$  holds in  $V$ ” are *not* equivalent.

- ▶ For **positive** extensions, they are equivalent for coherent implications.
  - e.g. the “Barr cover”.
- ▶ For **positive overt** extensions, they are equivalent for bounded first-order formulas.
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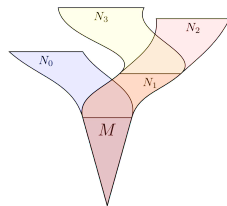
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- ▶ **somewhere** ( $\Diamond \varphi$ ) iff it holds in some positive extension.
- ▶ **proximally** ( $\Diamond \varphi$ ) iff it holds in some positive overt extension.

Foreshadowed by:

- 1984 André Joyal, Miles Tierney. *An extension of the Galois theory of Grothendieck.*
- 1987 Andreas Blass. *Well-ordering and induction in intuitionistic logic and topoi.*
- 2010s Milly Maietti, Steve Vickers. Ongoing work on arithmetic universes.
- 2011 Joel David Hamkins. *The set-theoretic multiverse.*
- 2013 Shawn Henry. *Classifying topoi and preservation of higher order logic by geometric morphisms.*





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For every inhabited set  $X$ , *proximally*  
there is an enumeration  $\mathbb{N} \twoheadrightarrow X$ .

For every ring, *proximally*  
there is a maximal ideal.

A preorder is well iff *everywhere*,  
every sequence is good.

A relation is well-founded iff *everywhere*,  
there is no descending chain.

A ring element is nilpotent iff  
all prime ideals *everywhere* contain it.

*Somewhere*,  
the law of excluded middle holds.

# The multiverse perspective

- 1 *Why has the inductive revolution been so powerful?*  
Because the inductive conditions are equivalent to truth in *all* forcing extensions.
- 2 *Why do proofs using Zorn's maximal ideals work so well in constructive algebra?*  
Because every ring proximally has a maximal ideal.
- 3 *Why are elements of  $\bigcap_p \mathfrak{p}$  not necessarily nilpotent?*  
Because we forgot the prime ideals in forcing extensions.
- 4 *How can we extract computational content from classical proofs?*  
By traveling the multiverse (upwards, keeping ties to the base), exploiting that
  - LEM holds *somewhere* and
  - DC holds *proximally*.

```

module _ (A : Set) where

open import Data.List
open import Data.List.Membership.Propositional
open import Data.Product

data Eventually (P : List A → Set) : List A → Set where
  now
    : {σ : List A}
    → P σ
    → Eventually P σ
  later
    : {σ : List A} {a : A}
    → ((τ : List A) → a ∈ (σ ++ τ) → Eventually P (σ ++ τ))
    → Eventually P σ

```

U:\*\*~ Countable.agda All L1 <N> (Agda:Checked +5 Undo-Tree)

U:%\*- \*All Done\* All L1 <M> (AgdaInfo Undo-Tree)

*Partial Agda formalization available.*